Hannenhalli-Pevzner Theory

- Signed, unichromosomal genomes
- Operation: reversal (signed)
- Polynomial time algorithms for distance & operations

• First polynomial result for a realistic model of genome rearrangements

- Genomes modeled as "permutations"
- What they call "permutations" are **not** functions from E to E
- Rather, they are functions from P (positions) to E (extemities)
- $P = \{1, 2, 3, ..., n\}$
- $E = \{1, 2, 3, ..., n, -1, -2, -3, ..., -n\}$
- π : P → E

- Reversals are permutations on P
- $\rho: P \rightarrow P$
- Reversals are applied to the **right**:
 πρ
- It is the only composition that makes sense
- Reversals do not always have small support

- Linear and circular cases are equivalent
- Extending the genome with 0 and n+1 essentially transforms the problem into a circular one
- In transforming π to σ , they fix σ :

 σ = 1 2 3 ... n or $\sigma(x) = x$ for all x

 σ is the "identity"

- The problem is then called "sorting" $\boldsymbol{\pi}$

• Formula:

 $d(\pi) = b(\pi) - c(\pi) + h(\pi) + f(\pi)$

- $h(\pi)$ = number of hurdles of π
- $f(\pi) = \begin{cases} 1 & \pi \text{ is a fortress} \\ 0 & \text{otherwise} \end{cases}$
- Algorithms O(n⁴) and O(n⁵), later improved

- Oriented pairs
- Score
- Algorithm 1: perform oriented reversals with maximum score as long as possible

- After Algorithm 1, one ends up with a "positive permutation"
- Reduced "permutations"
- Framed intervals
- Hurdles: cutting and merging

- Algorithm 2:
- 2k hurdles:
 - merge two hurdles, nonconsecutive if possible
- 2k + 1 hurdles:
 - simple hurdle:
 - cut it
 - no simple hurdle:
 - merge two hurdles, nonconsecutive if possible

- Algorithm 2, simplified:
- 2k + 1 hurdles and simple hurdle:
 - cut it
- Else:

- merge two hurdles, nonconsecutive if possible

- Final algorithm:
- while π is unsorted do
 - Algorithm 1
 - Reduce
 - Algorithm 2