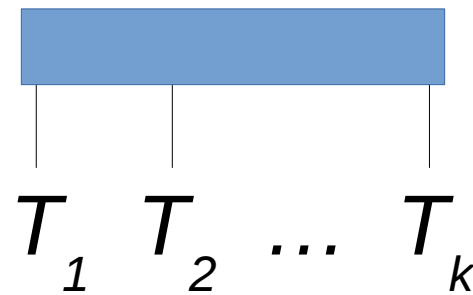
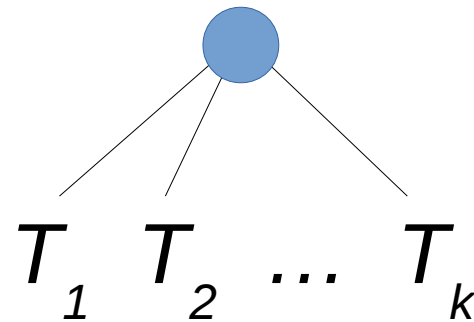


Testing for C1P using PQ-Trees

- C1P: order U so that certain sets $S \subset U$ are consecutive
- PQ-Tree: represents admissible permutations
- Online algorithm: tree changes with every new set added
- Reduction algorithm: linear in $|S|$ under amortized analysis
- Size of input: $n + m + f$

PQ-Trees

- Universal set U
- Leaves: elements of U
- Internal nodes:
 - P nodes
 - Q nodes



Proper PQ-Trees

- Every element of U appears exactly once as a leaf
- Every P node has at least 2 children
- Every Q Node has at least 3 children

Frontier

- Reading the leaves from left to right

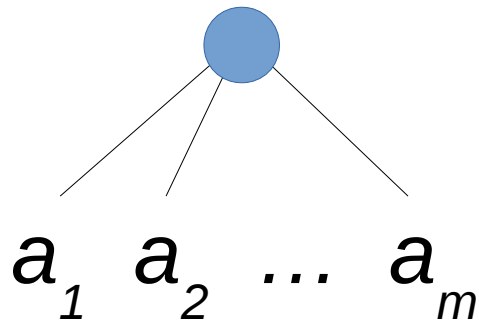
Equivalent trees

- Notation: $T_1 = T_2$
- T_1 and T_2 are equivalent when one can be transformed into the other by zero or more equivalence transformations:
 - arbitrarily permute the children of a P node
 - reverse the children of a Q node

$$\text{CONSISTENT}(T) = \{\text{FRONTIER}(T') \mid T' = T\}$$

Universal tree and null tree

- Universal tree: $|U| = m$



- Null tree = \emptyset $\text{CONSISTENT}(\emptyset) = \emptyset$

Modern implementation

PQTree<E>

universalSet: set of <E>
root: PQNode<E>

proper(): boolean
frontier(): list of <E>
equivalent(PQTree<E>):
 boolean
consistent(): set of
 <list of <E>>
isUniversal(): boolean
reduce(set of <E>)

PQNode<E>

parent: PQNode<E>
childList: list of
 <PQNode>
type: Enum(P,Q,L)

Reducing PQ-trees

- Input: PQ-tree T , set $S \subseteq U$
- Bottom-up: starts at leaves; process each node only after all its children have been processed
- Processing each node:
 - Find a pattern that applies to it
 - Replace the pattern by the replacement
- Stop when:
 - No pattern for a node: return null tree
 - After processing a node that “contains” S

Modern implementation (2)

PQTree<E>

universalSet: set of <E>
root: PQNode<E>

proper(): boolean
frontier(): list of <E>
equivalent(PQTree<E>):
 boolean
consistent(): set of
 <list of <E>>
isUniversal(): boolean
reduce(set of <E>)

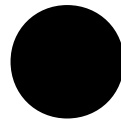
PQNode<E>

parent: PQNode<E>
childList: list of
 <PQNode>
type: Enum(P,Q,L)

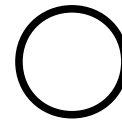
findPattern(): Pattern
apply(Pattern)
contains(set of <E>):
 boolean

Processing a node with respect to S

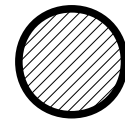
- Full, empty, partial



full



empty



partial

- Pertinent: full or partial
- Pertinent subtree: contains S
- $\text{ROOT}(T,S)$: lowest node whose frontier contains S

Modern implementation (3)

PQTree<E>

universalSet: set of <E>
root: PQNode<E>

proper(): boolean
frontier(): list of <E>
equivalent(PQTree<E>):
 boolean
consistent(): set of
 <list of <E>>
isUniversal(): boolean
reduce(set of <E>)
root(set of <E>): PQNode<E>

PQNode<E>

parent: PQNode<E>
childList: list of
 <PQNode>
type: Enum(P,Q,L)

findPattern(): Pattern
apply(Pattern)
contains(set of <E>):
 boolean
frontier(): list of <E>

Templates: pattern and replacement

- Goal: after replacement, the frontier of the subtree rooted at this node, as well as all equivalent subtrees, have all their pertinent leaves consecutive
- Affects node and its children only
- Type of node and label of children (full, empty, partial) matter
- Types of children do not matter



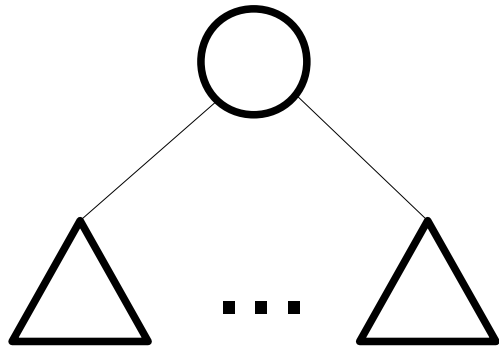
any type

Templates for leaves

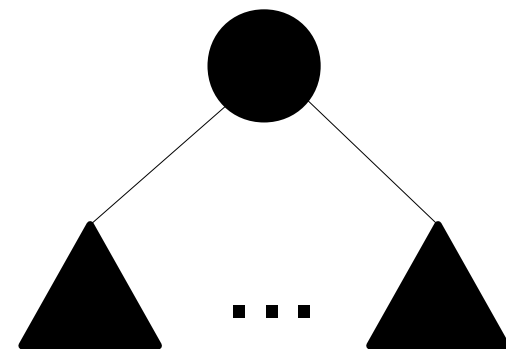
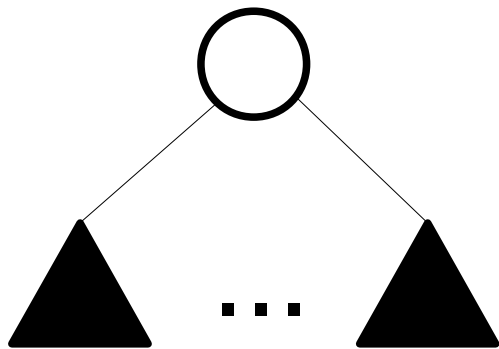
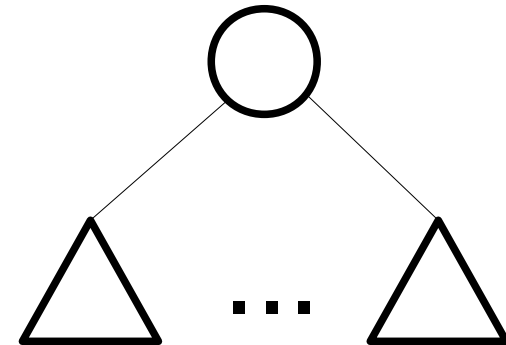
- Just label as empty or full

Templates for P-nodes

Pattern



Replacement



Templates for P-nodes (cont.)

- Templates when some children are empty and some full
 - Depend on node being $\text{ROOT}(T,S)$ or not
 - If node is not $\text{ROOT}(T,S)$, creates a partial P-node
- Templates when there are partial children
 - One or two partial children ok
 - More than two partial children: no pattern
- All templates: apply to equivalent nodes as well

Templates for Q-nodes

- Templates for all children empty or all children full
- Templates for some children empty and some full
- Depend on node being $\text{ROOT}(T,S)$ or not
- Templates for partial children
 - One or two partial children ok
 - More than two partial children: no pattern
- All templates: apply to equivalent nodes as well

Efficient implementation

- Need to avoid empty subtrees to guarantee amortized $|S|$ time
- Pruned pertinent subtree
- Two bottom-up phases
 - Bubbling up: identifying pruned pertinent nodes
 - Reduce: reducing pruned pertinent nodes
- Problem with parent pointers: blocking
- NORM(T): potential for amortized analysis