

# Solving C1P components

- Connected components of Strictly Overlap Graph
- Process each row in turn
- Place each row according to a previously placed neighbor
- To fix direction, need to have already processed neighbor of neighbor
- Process rows in DFS

# Algorithm *Place*

- **input:**  $w \rightarrow v \rightarrow u$  ( $w$  and  $v$  can be nil)
- **if**  $v = \text{nil}$  **and**  $w = \text{nil}$ : place 1s of  $u$  consecutively
- **if**  $w = \text{nil}$ : place  $v \setminus u, v \cap u, u \setminus v$
- **if**  $|u \cap w| < \min(|u \cap v|, |v \cap w|)$ :  
     $w, v \setminus u, v \cap u, u \setminus v$
- **else:**  
     $v \setminus u, v \cap u, u \setminus v, w$
- Check consistency of columns sets

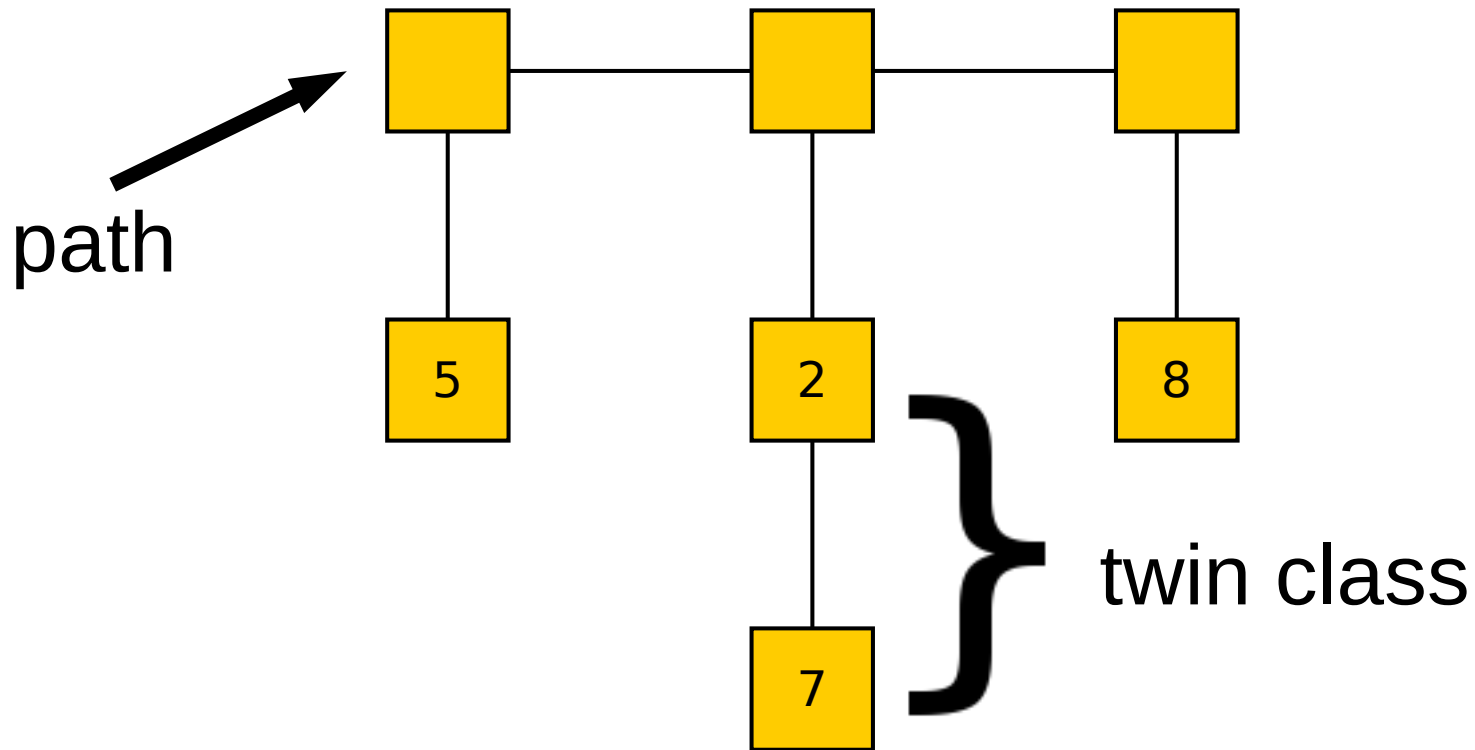
# Issues

- Twin elements: belong to exactly the same sets
- Twin elements are grouped in twin classes (column sets)
- How to handle twin sets?
- Complexity:  $O(mn) \rightarrow O(m + n + f)$ , where

$$f = \sum |S|$$

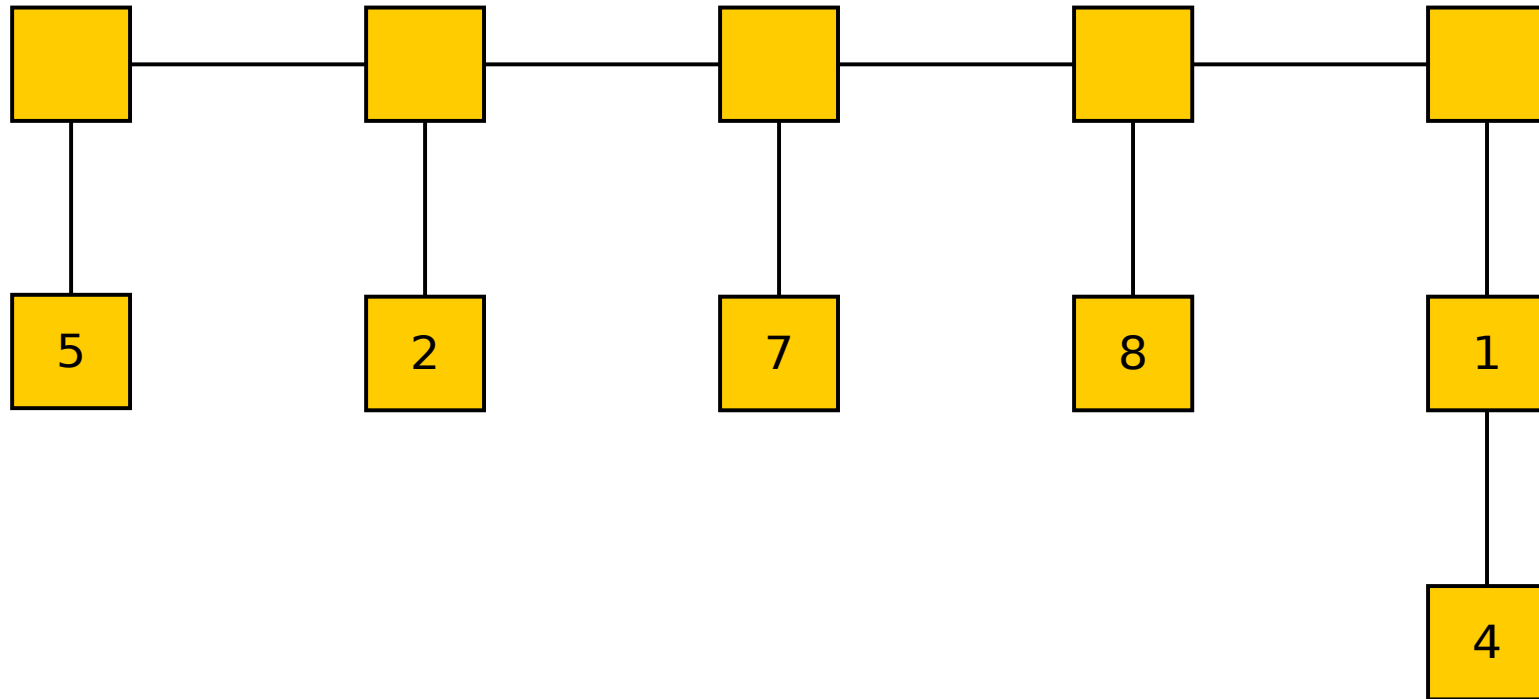
# Twin classes in a Path

- After adding  $\{2,7,8\}$  and  $\{2,5,7\}$



# Twin classes in a Path

- After adding  $\{2,7,8\}$  ,  $\{2,5,7\}$  , and  $\{1,4,7,8\}$



# Improved placement algorithm

- $\text{path} \leftarrow []$ ;       $\text{C1P} \leftarrow \text{true}$
- **for each** set  $S$  in DFS order :
  - color each  $x \in S$  in its class (possibly newclass)
  - for each** partial class  $C$  :
    - cut uncolored elements from  $C$  make a new class
    - place uncolored subclass near empty neighbor
    - if** no such neighbor :  $\text{C1P} \leftarrow \text{false}$
  - if** nonconsecutive full classes :  $\text{C1P} \leftarrow \text{false}$
  - if** there is a newclass :
    - if** both path extremities empty :  $\text{C1P} \leftarrow \text{false}$
    - place newclass in path extremity, preferably full
  - remove colors